# Minimization-Maximization Problems: Applications (in Communication), Challenges and Algorithms

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#### Outline

1 Min-max problems and motivation

2 The proposed solutions

- 3 Theoretical guarantees
  - 4 Numerical Results

## Mini-max problems

This talk mainly focuses on the following optimization problem

 $\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} f(x;y)$ 

f is some (possibly complicated) function over  $\boldsymbol{x},\,\boldsymbol{y}$ 

- x is the usual opt variable, power, precoder design, etc;
- $\min_x f(x,\cdot)$  is the usual cost minimization, i.e.,

cost = -throughput, or delay, *etc*.

• y is used to model provisioning of fairness, robustness, resilience, etc.

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#### Mini-Max Problems



Figure: Left convex/concave min-max problem; Right: Non-convex/Concave min-max problem

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#### Motivation from SP/Comm. perspective

- Question: Why Mini-Max problems?
- The "min" optimizes system level performance; while the "max" provides support such as fairness, robustness, resilience
- "min-max" together helps understand the scenario where some "adversary" (jammer) exists
- Lots of recent interests in this problem, applications in wireless transceiver design, adversarial (GAN)/robust learning, etc.

- Setting: MIMO interference channel with K users
- <u>Goal</u>: Design beamformers to maximize the min-rate utility under power and outage constraints



Figure: A set of transmitter-receiver pairs over an interference channel.

• The problem is given by  $(\mathbf{x}_i \text{ is user } i \text{ 's transmitter/receiver})$ 

$$\max_{\mathbf{x}} \min_{i \in [K]} R_i(\mathbf{x}_1, \cdots, \mathbf{x}_K)$$

•  $R_i(\mathbf{x})$  can be highly non-convex:  $\mathbf{X}_i$ : the transmit covariance matrix

$$R_{i}(\mathbf{X}) = \log \det \left( \mathbf{H}_{ii} \mathbf{X}_{i} \mathbf{H}_{ii}^{H} \left( \mathbf{I}_{N_{r}} + \sum_{l \neq i} \mathbf{H}_{li} \mathbf{X}_{l} \mathbf{H}_{li}^{H} \right)^{-1} + \mathbf{I}_{N_{r}} \right)$$

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• Connection to mini-max problem we just need to:

- Flip the sign:  $\min_{\mathbf{x}} \max_{i \in [K]} (-R_i(\mathbf{x}))$
- $\bullet\,$  Add a variable  ${\bf y}$  lives in a simplex, and equivalent formulation

$$\min_{\mathbf{x}} \max_{\mathbf{y}} \quad -\sum_{i=1}^{N} R_{i}(\mathbf{x}) y_{i} := -\mathbf{y}^{T} \mathbf{R}(\mathbf{x})$$
  
s.t. 
$$\sum_{i} y_{i} = 1, \quad y_{i} \ge 0, \ \forall \ i$$

- Problems with known global optimality
  - Max-Min SNR optimization [Zander, 1992] [Foschini and Gans, 1998]
  - MISO BF [Wiesel et al., 2005][Bengtsson and Ottersten, 1999]
  - Joint downlink BS association and power control [Sun-Hong-Luo 14]
  - Many more

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  - Many more
- More recent works involving non-convexity (so, no global min)
  - MISO coordinated beamforming with outrage constraints [Li et al., 2015]
  - MIMO coordinate transceiver design [Liu et al., 2011] (single stream) [Razaviyayn et al., 2011] (multi-stream)
  - MIMO constant envelop transceiver design [Shao et al 19]
  - Many more...

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- For problems that are global solvable, the standard approach: solve (a sequence of) convex problems like SDP/SOCP
- For non-convex problems, two popular ways in literature
  - Approximate the mini-max objective
  - 2 Translate to an "envelop form"

#### Popular Approximation for Non-Convex Min-Max

• One way to simplify is to use the log-sum approximation

$$\max_{i} -R_{i} \approx \frac{1}{\gamma} \log_2 \left( \sum_{i=1}^{K} 2^{-\gamma R_i} \right)$$

- A smooth approximation, large  $\gamma$ , good approximation
- Performance degradation

#### Popular Approximation for Non-Convex Min-Max

 The other way is to introduce an equivalent "envelop form" [Razaviyayn et al., 2011]

$$\min_{\lambda,\mathbf{x}} \ \boldsymbol{\lambda}, \quad \text{ s.t. } \quad -R_i(\mathbf{x}) \leq \boldsymbol{\lambda}, \ \forall \ i, \ \mathbf{x} \in X$$

- Reduces to a minimization form
- But still challenging, involving multiple non-convex constraints

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#### Example 2: Communication in the Presence of Adversary

- How to understand system performance/dynamics of communication systems with adversary?
- Regular user: minimize the cost; Jammer: maximize the cost

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- How to understand system performance/dynamics of communication systems with adversary?
- Regular user: minimize the cost; Jammer: maximize the cost
- An example [Gohary et al., 2009]: Interfering channel, N parallel tones, K users, optimization variables x<sub>k</sub><sup>n</sup>'s (power allocation of users over the tones), 1 jammer, optimizes y<sup>n</sup>

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} \sum_{(k,n)} -\log\left(1 + \frac{h_{kk}^n x_k^n}{\sigma^2 + \sum_{j=1, j \neq k}^K h_{jk}^n x_j^n + h_{0k}^n y^n}\right)$$

# Example 3: Robust Learning (over multiple domains)

• Empirical (non-convex) risk minimization (for training ML models)

$$\min_{\mathbf{x}} \frac{1}{N} \sum_{i=1}^{N} f_i(\mathbf{x})$$

- Treating all data equally/similarly
- In practice, the same model  $\mathbf{x}$  is used for multiple domains
- A model for digits recognition can be used in
  - identify the handwritten digits
  - ecognize the printed digits (e.g., house number)
- A "robust" model has to deal with both of domains [Qian et al., 2018]

### Example 3: Robust Learning (over multiple domains)

- Let the data draw from qth domain has lost function  $\{f_i^q(\mathbf{x})\}_{i=1}^N$
- When there are Q sets of data drawn from different domain (to describe the same phenomenon) [Qian et al., 2018]

$$\min_{\mathbf{x}} \max_{\mathbf{y}} \sum_{q=1}^{Q} \left( \sum_{i=1}^{N} f_i^q(\mathbf{x}) \right) \mathbf{y}_q, \quad \sum_{q=1}^{Q} y_q = 1, \ y_i \ge 0$$

- $\bullet$  Here  ${\bf y}$  can be interpreted as an adversarial distribution
- Identifying the "importance" of different data sets

#### Example 4: Distributed Learning

• Again the training problem, but with  $\boldsymbol{K}$  distributed agents

$$\min_{\mathbf{x}} \frac{1}{K} \sum_{i=1}^{K} f_i(\mathbf{x}_i), \text{ s.t. } \mathbf{x}_i = \mathbf{x}_j, \text{ if } (i,j) \text{ neighbors}$$

#### Example 4: Distributed Learning

 $\bullet$  Again the training problem, but with K distributed agents

$$\min_{\mathbf{x}} \frac{1}{K} \sum_{i=1}^{K} f_i(\mathbf{x}_i), \text{ s.t. } \mathbf{x}_i = \mathbf{x}_j, \text{ if } (i, j) \text{ neighbors}$$

• Stacking all variables  $\mathbf{x} := [\mathbf{x}_i, \cdots, \mathbf{x}_K]$ , re-write the above problem

$$\min_{\mathbf{x}} f(\mathbf{x}), \quad \text{s.t.} \quad \mathbf{A}\mathbf{x} = 0$$

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where A is the network incidence matrix (neighboring relations)

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where A is the network incidence matrix (neighboring relations)

 $\bullet$  Introducing the dual variable  $\mathbf y,$  we have

$$\min_{\mathbf{x}} \max_{\mathbf{y}} \quad f(\mathbf{x}) + \mathbf{y}^T \mathbf{A} \mathbf{x}$$

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Compared with the min-only problem, mini-max problem is challenging:

- Competing objectives, how to measure the progress of algorithm?
- e How to characterize the solution quality (when involving non-convexity)?
- In communication/signal processing applications, we also want computationally efficient algorithms
- Can we extend the existing algorithms for minimization (like gradient descent) problem to this setting?

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It illustrate some of the challenges, consider the simple setting

$$\min_{\mathbf{x}} \max_{\mathbf{y}} \mathbf{y}^T R(\mathbf{x}), \text{ s.t. } \mathbf{x} \in X, \mathbf{y} \in \Delta$$

Suppose that a good algorithm for the min-problem available (e.g., WMMSE [Shi et al., 2011], FP [Shen-Yu 18], Pricing [Shi et al 08])

$$\min_{\mathbf{x}} \mathbf{y}^T R(\mathbf{x}), \text{ s.t. } \mathbf{x} \in X$$

A natural approach: Alternatingly perform

$$\mathbf{x}^{r+1} = \mathsf{Algorithm-Step}(\mathbf{x}^r, \mathbf{y}^r), \quad \mathbf{y}^{r+1} = \max_{\mathbf{y} \in \Delta} \mathbf{y}^T R(\mathbf{x}^{r+1})$$

- Unfortunately, does not work
- The max step only selects a single user, sets everyone else to zeros
- How about making y step less greedy (perform one gradient ascent)?

 $\mathbf{x}^{r+1} = \text{Algorithm-Step}(\mathbf{x}^r, \mathbf{y}^r), \quad \mathbf{y}^{r+1} = \left[\mathbf{y}^r + \gamma R(\mathbf{x}^{r+1})\right]^+$ 

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Figure: Gradient Descent/Ascent dynamics exhibit oscillations (for any stepsize)

The classical literature in optimization has been concentrated in the case where f is convex/concave

- Extragradient method for finding saddle points in control [P. Vasilyev et al., 2010]
- Subgradient method for saddle-point problems [Nedić and Ozdaglar, 2009]
- Prox-method for smooth convex/concave problems [Nemirovski, 2005]
- Optimistic Gradient Descent Ascent [Daskalakis et al., 2017]
  - f bilinear, i.e  $f(x, y) = x^T A y$ ,
  - Gradient descent/ascent-type alg.  $\begin{aligned} x^{r+1} &= x^r - 2\alpha \nabla_x f(x^r, y^r) + \alpha \nabla_x f(x^{r-1}, y^{r-1}) \\ y^{r+1} &= x^r + 2\alpha \nabla_y f(x^r, y^r) - \alpha \nabla_y f(x^{r-1}, y^{r-1}) \end{aligned}$
- A few recent related works on non-convex/concave setting [Nouiehed et al., 2019], [Rafique et al., 2018], [Sanjabi et al., 2018a], [Sanjabi et al., 2018b]

#### Outline



2 The proposed solutions

- 3 Theoretical guarantees
  - 4 Numerical Results

• Intuitively, the "min" and "max" problems are not created equal

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- Intuitively, the "min" and "max" problems are not created equal
- Min outside, Max inside, we should allow the Max be solve relatively well before performing Min
- The alternating GD fails because the two sides are "equally powerful"
- The "exact-max" fails because one side that is "too powerful"



- A good algorithm should carefully balance between the two problems
- The min problem has to be slower than the max problem
- The max problem cannot be too aggressive
- Gradually adding regularizers to the min and max problems to control the speed of the two steps?

• Solution concept? First-Order optimality

$$\nabla_{\mathbf{x}} f(\mathbf{x}^*, \mathbf{y}^*) = 0, \ \nabla_{\mathbf{y}} f(\mathbf{x}^*, \mathbf{y}^*) = 0$$

or similar concepts to deal with constraints/non-smooth regularizers

• Call  $\epsilon$ -stationary solution if

$$\|\nabla_{\mathbf{x}} f(\mathbf{x}^*, \mathbf{y}^*)\| \le \epsilon, \ \|\nabla_{\mathbf{y}} f(\mathbf{x}^*, \mathbf{y}^*)\| \le \epsilon$$

 $\bullet$  Second-Order optimality:  $\mathbf{x}^*$  a "local min" and  $\mathbf{y}^*$  a "local max"



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### Min-max optimization problems

We consider a slightly more general min-max problem that involves

- Image: multiple blocks for the min problem
- e better modeling scenarios e.g., precoder design for K users; each problem has simpler structure

$$\min_{\{x_i \in \mathcal{X}_i\}} \max_{y \in \mathcal{Y}} \quad f(x_1, x_2, \cdots, x_K; y)$$

• assume that f smooth; non-convex w.r.t. x - concave w.r.t. y

# Hybrid Block Successive Approximation Alg. (HiBSA)

- A Hybrid block successive approximation algorithm
- Each time pick one block variable to perform descent/ascent
- Balance two regularization terms with coefficients  $\beta^r$  and  $\gamma^r$
- A simplified version to illustrate ideas
- 1 Perform K proximal gradient steps (Descent steps)

$$\arg\min_{x_i \in \mathcal{X}_i} \langle \nabla_{x_i} f(x_i^r, w_i^{r+1}, y^r), x_i - x_i^r \rangle + \frac{1}{\beta^r} \|x_i - x_i^r\|^2$$

or equivalently

$$x_i^{r+1} = \operatorname{proj}_{X_i} \left[ x_i^r - \beta^r \nabla_{x_i} f(x_i^r, w_i^{r+1}, y^r) \right]$$

2 Perform regularized ascent step (Ascent step)  $y^{r+1} = \arg \max_{y \in \mathcal{Y}} f(x^{r+1}, y) - \gamma^r ||y||^2 - \frac{1}{2\rho} ||y - y^r||^2$ 

# Hybrid Block Successive Approximation Alg. (HiBSA)

- Parameter choices  $\beta^r$  and  $\gamma^r$  are both diminishing sequences
- Intuition, the min steps slows down, allowing max problem to performs more steps per min update
- The max problem has some large regularization at the beginning, avoiding being too greedy at the beginning

# Hybrid Block Successive Approximation Alg. (HiBSA)

Now if we apply to the previous problem  $\min \max \mathbf{x}^T \mathbf{y}$ 



#### HiBSA - Extensions

- Can perform multiple ascent type steps in the maximization problem
- Do not need to perform gradient steps, but can solve some approximated minimization/maximization problems
- These extensions allows for flexible algorithm design, can plug and play existing minimization algorithms for the min-step
- For example, for min fair rate optimization algorithms, allow interlacing between WMMSE [Shi et al., 2011], or FP [Shen-Yu-18] steps, with ascent steps

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#### Outline



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# Convergence Guarantees of HiBSA

#### Assumptions

(1) For  $\{\gamma^r\}$  a diminishing sequence:

$$\gamma^r \to 0, \quad \sum_{r=1}^{\infty} (\gamma^r)^2 = \infty$$

(2) For  $\{\beta^r\}$  also a diminishing sequence, but diminishes faster then  $\gamma$ :

$$\beta^r = \mathcal{O}\left((\gamma^r)^2\right)$$

A typical choice:  $\gamma^r = \frac{1}{r^{1/4}}$ ,  $\beta^r = 1/\sqrt{r}$ 

#### Convergence Guarantees of HiBSA

#### Theorem (Convergence of HiBSA - Concave case)

For a given  $\epsilon > 0$  let  $T(\epsilon)$  be the minimum number of iterations needed to reach an  $\epsilon$ -stationary solution. Then we have

$$\boldsymbol{\epsilon} = \mathcal{O}\left(\frac{\log(T(\epsilon))}{\sqrt{T(\epsilon)}}\right)$$

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#### Section 4

#### **Numerical Results**

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#### Max-min fairness coordinated beamforming design

- Setting : MISO interference channel [Li et al., 2015]
- <u>Goal</u> : Design beamformers in order to maximize the min-rate utility under power and outage constraints



Figure: A set of transmitter-receiver pairs over an interference channel.

#### Max-min fairness coordinated beamforming design

#### Formulation :

 $\max_{x_i \in \mathbb{C}^{N_t}, \forall i} \min_i R_i(\{x_k\}) \quad \text{s.t. } \|x_i\|^2 \leq \bar{p}_i, \ \forall i + \text{ outage prob. constr.}$ 

- $x_i \in \mathbb{C}^{N_t}$  beamforming vectors,  $N_t$  no of transmitter antennas
- $\bar{p}_i$  power constraints
- <u>Solution</u> [Li et al., 2015] :
  - Adopt a suitable surrogate function for the utility function
  - Substitute the inner min problem with the log-sum-exp approximation,

i.e 
$$\min_{i} R_{i} \approx -\frac{1}{\gamma} log_{2} \left( \sum_{i=1}^{K} 2^{-\gamma R_{i}} \right)$$

• Solve the resulting problem exactly using CVX.

#### Max-min fairness coordinated beamforming design

Approximate problem; solve exactly with CVX [Li et al., 2015]
HiBSA [Proposed method]



Figure: Min-rate utility and runtime w.r.t noise level

Proposed method achieves higher min-rate utility and is significantly faster.

#### Max-min fairness linear transceiver design

- Setting : MIMO interfering broadcast channel in a multicell cellular network [Razaviyayn et al., 2011]
- <u>Goal</u> : Design beamformers in order to maximize the min-rate utility under power constraints



Figure: The interfering broadcast channel model

#### Max-min fairness linear transceiver design

• Formulation (envelop formulation) :

$$\min_{U,V,W,\lambda} \lambda$$
  
s.t.  $Tr[W_{i_k}E_{i_k}] - \log \det(W_{i_k}) - d_{i_k} \leq -\lambda, \ \forall i_k \in \mathcal{K}$   
$$\sum_{i \in \mathcal{I}_k} Tr[V_{i_k}V_{i_k}^H] \leq P_k, \forall k \in \mathcal{K}$$

- $i_k$  : *i*th user in cell k;  $\mathcal{K}$  : set of all cells;  $\mathcal{I}_k$  : set of users in cell k
- $E_{i_k}$  : MSE for user  $i_k$ ;  $d_{i_k}$  : no of data streams to  $i_k$ ;  $P_k$  : power const.
- V, U : transmit/receive beamformers;  $W, \lambda$  : auxillary variables
- Rate of user  $i_k = \max_{U_{i_k}, W_{i_k}} \log \det(W_{i_k}) Tr[W_{i_k}E_{i_k}] + d_{i_k}$
- Solution [Razaviyayn et al., 2011] :
  - $\bullet\,$  Solve the U,W subproblems exactly utilizing closed-form solutions
  - $\bullet\,$  Solve the V subproblem exactly using CVX.

Note that previously we resorted to an approximation, whereas here webace

#### Max-min fairness linear transceiver design - Results

- Envelop based solution [Razaviyayn et al., 2011]
- HiBSA [Proposed method]



Figure: Min-rate and runtime w.r.t number of cells/base stations

Proposed method achieves comparable utility and is clearly faster.

#### Power Control in the Presence of Jammer

• Setting : Parallel interference channel model [Gohary et al., 2009]

• <u>Goal</u> :

- Users : Maximize their individual rates
- Jammer : Reduce the total sum-rate of the other users
- Methods :
  - Interference pricing method (setting w/o jammer)
  - WMMSE algorithm (setting w/o jammer)
  - Regularized gradient descent/ascent (with jammer) [Proposed method]

#### Power Control in the Presence of Jammer-Results



Figure: Sum-rate w.r.t number of channels (left) and number of iterations (right).

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#### Conclusion

- Mini-Max is an interesting optimization problem arises in many contemporary applications
- SP/Comm problems, learning problems, GAN
- More challenging to analyze than the min-only problems
- Preliminary step towards understanding efficient algorithms; other related recent works [Nouiehed et al., 2019], [Rafique et al., 2018], [Sanjabi et al., 2018a], [Sanjabi et al., 2018b]
- Many open problems both sides non-convex? how to characterize solution quality, etc.
- Our paper can be found online arXiv preprint arXiv:1902.08294

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# **Thank You!**

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